The Use and Performance of Hashing Algorithms Sean Patrick Sanders 325G Jacobs Management Center SUNY-Buffalo Buffalo New York 14260 spsander@gmail.com

#### The Use and Performance of Hashing Algorithms<sup>1</sup>

#### Abstract:

Hashing concepts are an important, and often difficult, part of teaching computer science. Secure hash algorithms, which are used to verify that data has not been altered via man-in-the-middle threats, is also utilized for password protection, digital signatures, and to verify currency transactions in distributed blockchain ledgers. It is important that students receive a solid foundation in the applications of hashing algorithms. This paper presents four online exercises that can be used to illustrate secure hashing concepts that were developed in PHP. Understanding the interrelationships of blockchain, digital currency and hashing concepts are often difficult. But these set of experiential exercises facilitate the understanding of these concepts. Focusing on blockchain technology is a strong motivating force for understanding hashing concepts because of its current popularity.

#### **Secure Hashing Applications**

Hashing concepts are an important and sometimes difficult part of teaching computer science. The family of secure hash algorithms (SHA) used in cryptography has a variety of overlapping uses including:

- Ensures that data on servers has not been changed. Hashing significantly reduces processing time. Hash comparisons replace character-by-character comparisons of files and text to make sure that data has not been compromised.
- **Password protection.** The hashes are stored instead of the original password. Even if the password files are compromised, the password is not easily recovered by the attacker.
- **Digital signatures and fingerprinting.** The validity of downloaded software can be checked against a hash on the software developer's website.
- **Proof of work in digital currency mining**. Transactions are hashed until a certain number of leading zeros are obtained.

Secure hashing algorithms, from a student's perspective, are relatively complex. The first part of the paper will present the important concepts used in the hashing process. The second part will describe a portfolio of applications that reinforce and explain the secure hashing concepts, and in particular, blockchain concepts. As noted earlier, focusing on blockchain technology is a strong motivating force for understanding hashing concepts because of its current popularity.

# **Blockchain Technology**

Blockchain technology is revolutionizing the way that individuals are handling and interacting with cryptocurrency. Many organizations and individuals are trying to get involved in mining, but the tremendous computational and energy demands of mining are reducing the opportunity for success. The decrease in mining success for new entrants is related to the presence of large

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mining companies and consortia with deep pockets. A mining pool is a group of miners that join forces to combine computing power for monetary gain (Figure 1). Large mining pools increase the chance for successful hashing.

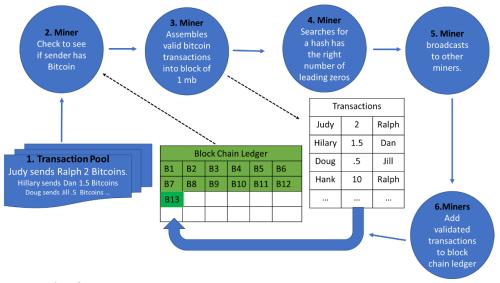


Figure 1: The Mining Process

The underlying hashing algorithm behind blockchain mining can be traced to the 1970s. Sophisticated hashing algorithms did not emerge until the late 1990s. They included the original SHA 1 and MD5 hash algorithms. These algorithms are essential in both computing and blockchain mining because they hold the key to security and authenticity. The authenticity of data is maintained by ensuring the data has not been changed; this is accomplished by using a complex mathematical algorithm. The hash algorithm is also used in mining to make sure that there is no cheating or double spending and to eliminate the presence of spammers. A spamming attack during the mining process will reduce the efficiency and overall hash rate of the mining process. The mining process is computationally intensive and involves the difficult task of creating a hash with leading zeros. Difficulty in mining increases as the number of leading zeros increases. At this point in time, Bitcoin miners are required to produce hashes with 17 leading zeros.

# **Dedicated Hardware Mining**

The computational requirements of Bitcoin mining are significant and are a major reason why blockchain implementations do not readily scale. Bitcoin implementations are very time consuming to setup and require powerful servers and extended infrastructures. Users must not only be concerned with the GPU processing power and RAM, but the power supply, motherboard, and graphics card as well. The computational intensiveness of Bitcoin mining has led to the use of Application-Specific Integrated Circuits (ASICS) that are expensive, loud, have a short life, and produce substantial power costs. It has been estimated that a bitcoin transaction consumes more than 5,000 times more energy than a credit card transaction.

Ethereum is in the process of using a different hash algorithm, <u>Ethash</u>, for proof of work; it is ASIC resistant and permits block verification by a light client. Eventually, GPUs may not be needed to mine digital currency. Rather a grid computing network could be employed involving millions of connected computers to engage in transaction verification.

### **Details on the Hashing Process**

A hash function is used to verify that data has not been changed. A hash function is used to map data back to a specific set of data of a predetermined size. If you want to check if any value in a stream of characters has changed, you can check the hash value. A verified hash value indicates that the original data remains untouched. If the hash value does not match the hash value from the original data, then data has been altered or tampered with in some way. This is in part how the double spending problem is countered, and man-in-the-middle attacks are mitigated. It should be noted that checksum computation has some similarities, but checksums are different because they are not unique, whereas the SHA values are essentially inimitable. And while hash collisions do occur, they are rare with recent implementations of hash algorithms. A simple example of how hashes are generated is shown in Figure 1.

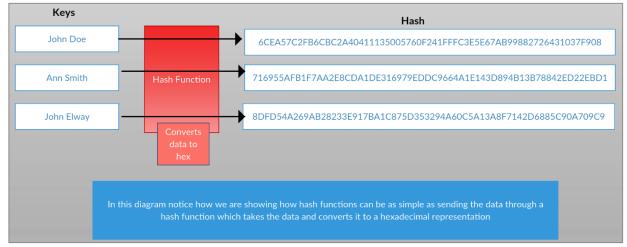
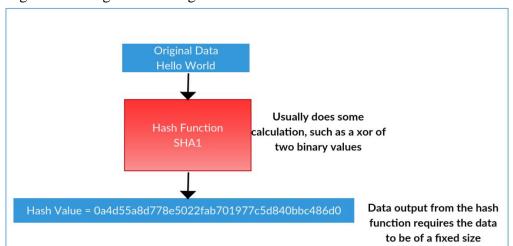


Figure 2: Hash Algorithm Example Developed by the Authors

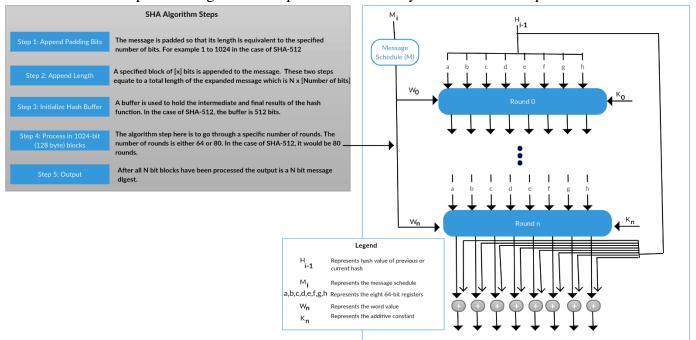
Hashing algorithms are utilized in computer security. In the case of the blockchain, they are used to ensure that the data blocks have not been modified and act as a verification system for miners. In computer security programs such as antivirus programs, hashes are used to make sure that the applications on your personal computer are not hacked or manipulated in some way. If the hash values of programs and files are different than what is contained in the antivirus database, they will be flagged. The antivirus program checks all installed applications to determine if the hashes in the database match the hashes computed during the virus check. There are many other



applications and software that use hashing algorithms to ensure the integrity of data. Below is a diagram showing how hashing works.

Figure 3: How Hashing Works

Hashing using SHA involves multiple steps, whether it be SHA 224, 256, or 512 algorithms. The first step in SHA is to append padding bits. This means adding a single 1 bit followed by the necessary number of zero bits. For example, if we have 1101 and we require padding of two additional bits, we would get 110100. The key is always to add the zeros padding at the end. The next step is to append the length, which is dependent on the block size. This means we have to use the following rule of  $0 \le k \le 512$ . The k in this instance should be equivalent to 448, and thus we conclude the formula can be  $448 \equiv -64 \pmod{512}$ . In any case, the 512 doesn't have to be 512 but can change, depending upon the use of SHA 224 or 256. This step is required to distinguish the empty input from the longer input. The next step is to initialize the hash buffer by representing eight 64-bit registers, each consisting of hexadecimal values (a,b,c,d,e,f,g,h). These values represent the 9th through 16th primes. The formula is as follows:  $0=|\operatorname{frac}(\sqrt{n}) \cdot * 264|$ . The letter n, in this case, represents the prime number. The next step is to process the message in nbit blocks. The n-bit depends on the block size. The number of rounds is either 64 or 80, depending on which SHA algorithm you choose to use. In each round, you take as an input the buffer value of the previous step and a sum of both the buffer value of the previous step and the previous hash value. The final output result in the last round step is when you have completed



each of the above steps. The diagram below provides a summary of all the SHA steps.

*Figure 4: How SHA Algorithm Works (Adapted and redrawn from Cryptography and Network Security Principles and Practice Seventh Edition)* 

#### Hashing Applications for Teaching

A set of four exercises were developed to illustrate hashing concepts. These exercises are coordinated by the instructor. However, class participation is the cornerstone of the exercises. Students should bring a laptop to class, but a smartphone with access to the internet will also suffice. The exercises all use PHP so that they can be run on virtually any device.

#### **Exercise 1: Generating a Nonce**

The purpose of the first exercise is to illustrate the output from a SHA256 hash and to show how a nonce is used. Nonces are used extensively in cryptographic authentication. In several blockchain implementations, the nonce is a unique value that is added to the end of the text that is being hashed with the objective of generating a hash with a specified number of leading zeros. Bitcoin currently requires about 17 leading zeros. The program used in these exercises is available at <a href="http://104.156.254.129/Exercise1.php">http://104.156.254.129/Exercise1.php</a>.

The first step in Exercise 1 is to tell the students to enter their name. This will generate the SHA256 hash. Then explain that it does not matter how many characters are entered; the hash size for each algorithm will be the same. A fixed size hash will be created. Also, explain that it is nearly impossible to take the hash and reverse engineer it to find the original text. This is particularly true of the newer SHA implementations, such as SHA256 and SHA384. You can also note that the SHA256 and SHA384 are more secure than the other algorithms.

The second step in Exercise 1 is to change the first letter of their last name to a lower case. Ask participants to compare the two hashes, which will not be the same.

The third step is to have students enter their name along with the number 1 right after their name. Explain that the number following their name is called a nonce. They will click on the submit button and refer to the SHA256 result. If there is no leading zero tell them to put a 2 after their name and see if there is a leading zero for the hash. This should be repeated until a SHA256 hash with a leading zero is generated. On average, this will take approximately 16 iterations. Figure 5 illustrates how the MD2 hash with a leading zero for Donna Summer with a nonce value of 29 was generated.

<u>Generating a Hash</u>					
	Donna Summer29				
Enter in Sentence:	G				
Submit					
	Value of: Donna Summer29 )da88f4c858946271269f188a5218e2e08e82e7c	7f2bfc7dc2d56bc6b5			

Figure 5: Nonce Generation

At this point, you can also have students copy a large amount of text from a web page or a file into the text box. Then they submit the text to the hash program. Tell them that the hash is still the same size for each algorithm. Ask them to change any letter in the text to another letter or number, and notice any change in the hashes.

Here is a summary of the steps in Exercise 1:

- 1. Enter your first and last name into the submit box and click Submit
- 2. Change the first letter in your last name from upper to lower case and click Submit.
- 3. Enter your first, and the last name followed by the number 1 into the text box and then click Submit. If the first character is a zero stop. Capture the screen and save it. If the first character was not a zero, keep incrementing the number following your name by one more unit until you generate a hash with one leading zero. Capture the screen and save it.
- 4. Cut and paste some data into the text box and click Submit.
- 5. Change any letter in the text and Submit.
- 6. Extra credit: keep adding numbers until you get two leading zeros. (Don't try anything beyond two leading zeros, for it will take a long time.)

In Bitcoin mining, the SHA256 hash is generated by adding a nonce or unique value to the end of the text that is being hashed. We added that number to the end of the name. The goal is to generate a SHA256 hash that starts with zeros. The nonce is the random number that is added to the end of the text being hashed until the desired number of leading zeros are generated. This adding of a random number, or nonce, to generate a hash with leading zeros is what mining is all about.

It takes more time to generate a hash with four zeroes in front than eight zeroes. A hash with one zero requires about 16<sup>1</sup> or 16 attempts. A hash with two zeros requires about 16<sup>2</sup> or 256 attempts, while a hash with three zeros requires about 16<sup>3</sup> or 4096 attempts. Right now Bitcoin miners have to generate hashes with 17 leading zeros 16<sup>1</sup> or 2.9514791e+20.

#### Exercise 2: A Hashing Program that Automatically Searches for a Nonce

This program searches for a nonce for a character string. This algorithm is quite complex because the program has to keep searching until it finds a hash with a leading zero. There are issues related to converting numbers into strings and checking for leading zeros. Dedicated ASICs hardware and GPUs have fine-tuned these operations and are difficult to out-perform. This program is also written in PHP and is available at <a href="http://104.156.254.129/Exercise2.html">http://104.156.254.129/Exercise2.html</a>. On small amounts of text, this program will have hash rates over 400,000, which are not anywhere near the trillions of hashes per second of <u>ASICs</u> processors. Also note that this program selects a random number of the first nonce and then increments it by 1. This is more accurate view of how the mining process is performed.

Here is a summary of the steps for this exercise:

- 1. Enter your name and the number of Bitcoins you want to give to a friend.
  - a. For example, Sean transfers 2 Bitcoins to Matt
- 2. Enter 5 for the number of leading zeros to generate.
- 3. Enter SHA512 for the hashing algorithm.
  - ✓ What was total number of attempts?
  - $\checkmark$  What was the expected number of attempts?
  - ✓ How long did it take to find the nonce that generated the correct number of leading zeros?
  - $\checkmark$  What was the hash rate?

The instructor should ask the class who had the smallest number of attempts and who had the greatest number. He or she should then write them on the board. Also, it would be worthwhile to inquire about the hash rate times and why they were different. The hash rate is related to the availability of virtual machine resources, but it is still interesting to observe. Tell the class that dedicated mining processors, such as the <u>ANTMINER S9</u>, generate trillions of hashes per second until they find a hash with 17 leading zeros when mining Bitcoin.

Figure 6 presents the inputs, and Figure 7 presents the result for "Sean transfers 2 Bitcoins to Matt," using SHA512 and searching for 5 leading zeros.

<u>Searching for a Nonce</u>
This application illustrates how digital currency mining is done using hashing as proof of work. You enter in the text to be hashed, the number of leading zeros and click on the hashing algorithm desired. The program will find the hash by adding a nonce, or random number, to the string until it generates a hash with the appropriate number of leading zeros. There is there is a limit of 30 cpu seconds for the computation, so stick to 5 or less leading zeros.
Please enter in a sentence and select the hash you want to use
Enter in Sentence: Enter in Number of Zeros to Check For: 5 Select a Hash: 544512 T
Submit
Go Back
Figure 6: Input
Results from Searching for a Nonce
Your Sentence is: Sean transfers 2 Bitcoins to Matt Your sentence length is: 33 Zeros to account for is: 5 Nonce value is: 586,305,402 The initial hash of the sentence without Nonce is: e3c293bc3f125609ee1f86dfd15744766d96623fbc8ef0cdd38b2f50babdf6f2eebace442b65aa79622c0e5763a572fcb32a26a46f338b5de8bb6ccf62b01d9c You picked <u>SHA-512</u> for the hash algorithm
Initial Hash Value with Nonce is: 61a0e5e58b3d269758b1e8ae6a105d83bb3fc2f270ae2f81ce7aff195bb8df1e0dc8bcbc979557235bf621cfe8a4103db89a4b4a473e4661fe75881fccab7f7d Nonce value is: 586,305,402
Poculte:
Hesuits. Hash Values for last occurence is: 0000089fcc4557cf7a8f7e8za56a2zf55328e814c05af56f6474e2cb39c731bd78b784140e717c0b38ee0ea61d63da2efa87b6dae50dce57444397bb6a78f774 Nonce value is: 587,621.883 Total Number of attempts is: 1,316,481 The Expected number of attempts is: 1,048,576 Start Time is: 0,001 Seconds

Figure 7: Output

#### **Exercise 3: Mining Simulation**

The purpose of this exercise is to illustrate in greater detail the computational demand that is required for using hashing for proof of work. It requires participants to enter the text to be hashed, along with the number of leading zeros, then to click on the hashing algorithm desired and the number of times to run the simulation. The program will find the hash by adding a nonce, or random number, to the string until it generates a hash with the appropriate number of leading zeros. This program is available at <u>http://104.156.254.129/Exercise3.html</u> and it can be run using a laptop or mobile device with internet access. The simulation can be run with SHA256, SHA512, SHA384, or the older SHA224.

To start the process, students should enter their name and major, check for four zeros, select SHA256 for the hash and have the simulation run for ten iterations. The input is illustrated in Figure 8, and the result of the simulation is illustrated in Figure 9.

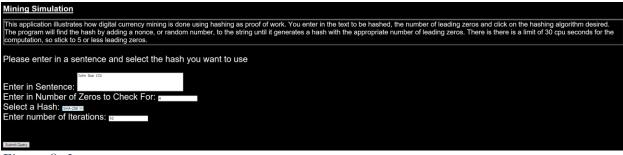


Figure 8: Input

Results from Mining Simulation Your Sentence is: John Doe CIS Your sentence length is: 12 Zeros to account for is: 4 Nonce value is: 933,449,540 The initial hash of the sentence without Nonce is: e4fcc0ca29e6c094daf393550cd44bca3f97d2aede7be9b3fb0511434748dcb5 You picked <u>SHA-256</u> for the hash algorithm								
Number of Attempts	Time in Seconds	Hash Rate						
70,422	0.16	440,137.500						
10,698	0.03	356,600.000						
21,017	0.05	420,340.000						
56,672	0.13	435,938.462						
60,458	0.14	431,842.857						
61,141	0.14	436,721.429						
93,395	0.22	424,522.727						
173,599	0.41	423,412.195						
90,083	0.21	428,966.667						
7,769	0.02	388,450.000						
Average Attempts	Average Time	Average HashRate						
64 525	0.151	418,693,180						

Figure 9: Output

Students can simulate the performance of the various algorithms and copy the results into a spreadsheet. These results can then be used to write a short paper that discusses the results. Figure 11 was constructed by cutting and pasting the results from running the simulation 100 times for all three hash functions, by varying the number of characters from 10 through 2,000. About 76,838,804 million hashes were used to generate these results. The projected number of hashes to compute was 78,643,200 (100 simulations x  $16^{A4}$  leading zeroes x 3 SHA algorithms x 4 different character lengths). This is a further illustration that the proof of work simulation is working correctly.

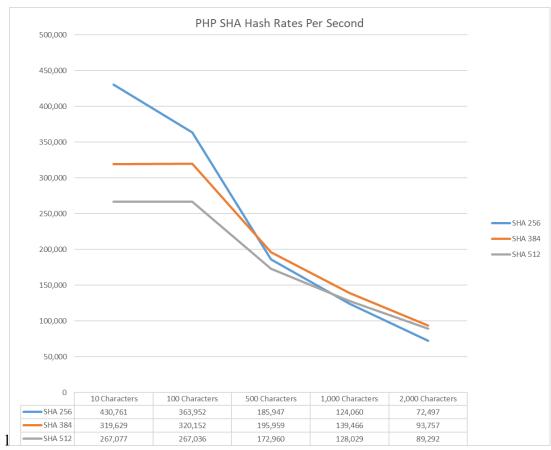


Figure 10: SHA Hashing Rates on a PHP server

There are many ways you can use this simulation. It can be used to understand how the number of leading zeros translates to computational intensity. The average number of attempts is a function of the number of leading zeros and is  $16^{n}$  where n is the number of leading zeros.

#### **Exercise 4: Birthday Paradox and Cracking Secure Hash Algorithms**

The safety of secure hash algorithms is always an issue of interest. The <u>Birthday Paradox</u> can be used as an approximation of the amount of brute force computing necessary to find a hash collision. The program used to illustrate the number of years to find a hash Collison for various SHA bit sizes and hash rates can be found at <u>http://104.156.254.129/BirthdayParadox.php</u>.

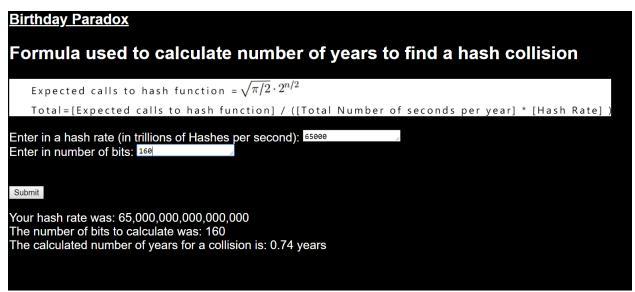
The birthday paradox is when you have a group of x people there exists a pair of people that have the same birthday. The birthday paradox can be applied to hash collision attacks in which two input strings can have the same hash result. The purpose of this exercise is to illustrate how the Birthday Paradox can help us determine the strength of the mining process. We can use the birthday paradox formula:

```
Expected calls to hash function = \sqrt{\pi/2} \cdot 2^{n/2}
Total=[Expected calls to hash function] / ([Total Number of seconds per year] * [Hash Rate] )
```

The formula above allows us to understand how long it would be before a hash collision occurs. For example, the number of years needed to find a collision for the lowly 160 bit SHA1

algorithm using 5,000 ASICS computers each capable of 13TH/s with a total hashing rate of 65,000 TH/s, is 0.74 years See Figure 10. This is in contrast to the 208.06 trillion years to find a collision with the SHA256 algorithm using the brute force approach.

Now let's consider just one S9 Antminer with a hash rate of 14 TH/s. It would take .05 years to find a collision using the S9 for the 128 bit SHA1 algorithm. It would take 9.66 trillion years to find a collision for the SHA256 algorithm. The students should be encouraged to play around with the number of bits and the hash rate. The purpose of increasing and decreasing the hash rate is to understand how the hash rate plays an important role in being able to understand how it affects the number of years for a collision to occur. Also, the decrease in the number of bits will then allow for a deeper understanding of why blockchain used SHA 256 over SHA 128 and eventually lead to stronger algorithms like SHA 3.





The 256, 384 and 512 bit hashing algorithms are considered secure at the moment. With the emergence of quantum computing, it may be possible to compromise these algorithms and break the current security mechanism underlying current blockchain implementations (E O Kiktenko, 2018). Any party that has access to a quantum computer will certainly have an unfair advantage of winning the mining the process and theoretically perform attacks that alter the digital signatures. Perpetrators could use a quantum computer generate a hash collision and alter the blocks data, and without any consequences. In essence, the original blockchain data could be altered, but there would be no alteration in the hash. This would ruin the security of the current blockchain infrastructure.

#### **Conclusion and Future Direction**

This project has explored the importance of secure hash algorithms in a variety of settings, and in particular, digital currency mining. Three exercises were developed to illustrate how secure hash algorithms are used to increase the security and integrity of data, to understand the various implementations of hashing algorithms, and to understand the mining process. A PowerPoint slide deck for teaching the three the exercises is available at <a href="http://104.156.254.129/">http://104.156.254.129/</a>

We are also developing a blockchain simulation called BARTS (Blockchain ART Simulation), where students can participate in simulations of a digital coin for buying and selling drawings. We have found that BARTS simulation should be introduced before presenting the material in this paper. Students need a conceptual image of the mining process before they can integrate the hashing concepts. The BARTS simulation will be used to illustrate market demand concepts, how transactions are processed on the blockchain, the role of smart contracts, how the ledger is updated, and how gas can be used to reduce transaction times. Figure 12 illustrates the basic mechanics of the BARTS process. We then introduce the hashing concepts discussed in this paper.

We have used the material discussed here to teach 130 Masters of Science MIS students hashing and blockchain concepts in three hours. At the end of the sessions, we had the students in the course fill out an anonymous survey on the modules. The survey questions and corresponding results are as follows.

**Question 1:** The material covered in the blockchain teaching module helped me gain a clearer understanding of blockchain concepts (1 Strongly Disagree to 5 Strongly Agree). The mean was 4.1 with 108 respondents

**Question 2:** What is the percentage of new understanding of blockchain concepts did you get from the teaching module?

N	1inimum %	Maximum %	Mean %	N
	17	100	68.18	108

The results of the survey were positive. These results are promising, given that many of the individuals in the class had extensive experience in systems development and applications programming.

Teaching hashing concepts is much easier when the teaching pedagogy is interesting. We think the approach used here establishes a strong foundation for understanding hashing concepts and also creates an opportunity for discussing contemporary blockchain concepts.

# **BARTS** Players

- Coordinator or instructor
- 3 Artists
- 3 Gallery Owners
- 3 Mining Pools with three miners in each pool. One of the miners will be the spokesperson for the mining pool.







2. The gallery owner takes the completed Bart's diagram to the coordinator.

The coordinator asks the mining pool to validate the Bart's transaction after determining if the gallery owner has sufficient number of Barts.

4. The mining pool hashes the transaction.

5. The mining pool that finds the hash with the correct number of leading zeros tells the coordinator they have found the hash.

The coordinator announces what nonce to use that will generate the hash with correct number of leading zeros.

7. All of the mining pools validate the transaction with the nonce.

8. If the mining pools validate the hash, then the Bart's sheet is put up on the front wall and the ledgers are all updated by the mining pools.











Figure 12: BARTS the Blockchain Art Simulation

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